Chapter 5 Structured Query Language

Content

- Introduction
- Data definition
- Data manipulation
 - Query
 - Update
- View definition
- Index

Introduction

- Relational algebra language
 - How to execute the query operations (in what order)
 - Difficult for users
- SQL (Structured Query Language)
 - High level declarative language interface
 - The user only specifies what the result is to be
 - Developed at IBM Research (1970s)
 - Also pronounced SEQUEL
 - Expanded to a standard by ANSI
 - SQL1: SQL-86
 - SQL2: SQL-92
 - SQL3: SQL-99, SQL-2003, SQL-2006, SQL-2008

Introduction

- SQL includes
 - Data definition
 - Data manipulation
 - View definition
 - Integrity constraint
 - Security and authorization
 - Transaction control
 - Rules for embedding SQL into programming languages

Content

Introduction

Data definition

- Data type
- Data definition commands
- Data manipulation
- View definition
- Index

Data definition

- Describes the structure of information in the DB
 - Schema for the relation
 - Domain of each attribute
 - Integrity constraint
 - Index on each relation
- Consists of
 - CREATE TABLE
 - DROP TABLE
 - ALTER TABLE
 - CREATE DOMAIN
 - CREATE DATABASE

^{- ...}

Data type

- Numeric
 - INTEGER
 - SMALLINT
 - NUMERIC, NUMERIC(p), NUMERIC(p,s)
 - DECIMAL, DECIMAL(p), DECIMAL(p,s)
 - REAL
 - DOUBLE PRECISION
 - FLOAT, FLOAT(p)

Data type

- Character string
 - CHARACTER, CHARACTER(n)
 - CHARACTER VARYING(x)
- Bit string
 - BIT, BIT(x)
 - BIT VARYING(x)

Datetime

- DATE
- TIME
- TIMESTAMP

Create table command

- Define a new relation by giving
 - A name
 - Attributes
 - Names
 - Data types
 - Integrity constraints on attributes
- Syntax

```
CREATE TABLE <Table_name> (
        <Column_name> <Data_type> [<Constraint>],
        <Column_name> <Data_type> [<Constraint>],
        ...
        [<Constraint>]
)
```

Example

CREATE TABLE EMPLOYEE (SSN CHAR(9), LNAME VARCHAR(10), MNAME VARCHAR(20), FNAME VARCHAR(10), BDATE DATETIME, ADDRESS VARCHAR(50), SEX CHAR(3), SALARY INT, SUPERSSN CHAR(9), **DNO INT**

Create table command

- Basic <Constraint>
 - NOT NULL
 - NULL
 - UNIQUE
 - DEFAULT
 - PRIMARY KEY
 - FOREIGN KEY / REFERENCES
 - CHECK

Give a name to constraints

CONSTRAINT <Constraint_name> <Constraint>

Example – Constraint

CREATE TABLE EMPLOYEE (LNAME VARCHAR(10) NOT NULL, MNAME VARCHAR(20) NOT NULL, FNAME VARCHAR(10) NOT NULL, SSN CHAR(9) PRIMARY KEY, **BDATE DATETIME**, ADDRESS VARCHAR(50), SEX CHAR(3) CHECK (SEX IN ('Nam', 'Nu')), SALARY INT DEFAULT (10000), SUPERSSN CHAR(9), DNO INT

Example – Constraint

CREATE TABLE DEPARTMENT (DNAME VARCHAR(20) UNIQUE, DNUMBER INT NOT NULL, MGRSSN CHAR(9), MGRSTARTDATE DATETIME DEFAULT (GETDATE()) CREATE TABLE WORKS_ON (SSN CHAR(9) FOREIGN KEY (SSN) **REFERENCES** EMPLOYEE(SSN), PNO INT REFERENCES PROJECT(PNumber), HOURS DECIMAL(3,1)

Example – Constraint name

CREATE TABLE EMPLOYEE (

```
LNAME VARCHAR(10) CONSTRAINT EM_LNAME_NN NOT NULL,
```

MNAME VARCHAR(20) NOT NULL,

FNAME VARCHAR(10) NOT NULL,

SSN CHAR(9) CONSTRAINT EM_SSN_PK PRIMARY KEY,

BDATE DATETIME,

ADDRESS VARCHAR(50),

```
SEX CHAR(3) CONSTRAINT EM_Sex_CHK
```

```
CHECK (SEX IN ( 'Nam', 'Nu' )),
```

SALARY INT CONSTRAINT EM_Salary_DF DEFAULT (10000),

```
SUPERSSN CHAR(9),
```

DNO INT

Example – Constraint name

CREATE TABLE WORKS_ON (

SSN CHAR(9),

PNO INT,

HOURS DECIMAL(3,1),

CONSTRAINT WO_SSN_PNo_PK PRIMARY KEY (SSN, PNO),

CONSTRAINT WO_SSN_FK FOREIGN KEY (SSN)

REFERENCES EMPLOYEE(SSN),

CONSTRAINT WO_PNo_FK FOREIGN KEY (PNO)

REFERENCES PROJECT(PNUMBER)

Alter table command

- Is used for modification
 - The structure of tables
 - Integrity constraints
- Columns

ALTER TABLE <Table_name> ADD COLUMN <Column_name> <Data_type> [<Constraint>]

ALTER TABLE <Table_name> **DROP COLUMN** <Column_name>

ALTER TABLE <Table_name> **ALTER COLUMN** <Column_name> <New_data_type>

Alter table command

Constraints

. . .

ALTER TABLE <Table_name> ADD CONSTRAINT <Constraint_name> <Constraint>,

CONSTRAINT <Constraint_name> <Constraint>,

ALTER TABLE <Table_name> **DROP** <Constraint_name>

Example

ALTER TABLE EMPLOYEE ADD JOBTITLE CHAR(20)

ALTER TABLE EMPLOYEE DROP COLUMN JOBTITLE

ALTER TABLE EMPLOYEE ALTER COLUMN JOBTITLE CHAR(50)

Example

CREATE TABLE DEPARTMENT (

DNAME VARCHAR(20),

DNUMBER INT NOT NULL,

MGRSSN CHAR(9),

MGRSTARTDATE DATETIME

)

ALTER TABLE DEPARTMENT ADD

CONSTRAINT DE_DNumber_PK PRIMARY KEY (DNUMBER),

CONSTRAINT DE_MgrSSN_FK FOREIGN KEY (MGRSSN)

REFERENCES EMPLOYEE(SSN),

CONSTRAINT DE_MgrStartDate_DF DEFAULT (GETDATE())

FOR (MGRSTARTDATE),

CONSTRAINT DE_DName_UN UNIQUE (DNAME)

Drop table command

- Is used for deleting the structure of tables
 - All the data in a table are also deleted
- Syntax

DROP TABLE <Table_name>

Example

DROP TABLE EMPLOYEE

DROP TABLE DEPARTMENT

DROP TABLE WORKS_ON

Example



Content

- Introduction
- Data definition

Data manipulation

- Basic queries
- Set, set/multiset comparison and nested queries
- Aggregate functions and grouping
- Others
- Data update
- View definition

Index

Query

- Data manipulation language is used for retrieving information from a database
 - These tuples often satisfy a certain condition
- Based on



- Allow a table to have two or more tuples that are identical in all their attribute values
- Not a set of tuples, but a multiset or bag

Basic query

Is formed of the three clauses

SELECT <List_of_columns>

FROM <List_of_tables>

[WHERE] <Condition>

- <List_of_columns>
 - Column names showed in the result of the query
- <List_of_tables>
 - Table names required to process the query
- <Condition>
 - Boolean expression that identifies the rows to be retrieved
 - Expression's connection : AND, OR, and NOT
 - Operations: < , > , <=, >=, <>, =, LIKE and BETWEEN

Basic query

SQL and Relational Algebra



SELECT L $FR\overline{O}_{C}(\mathbb{R})$ WHERE C

Example

The entire tuple is produced

SELECT *-----

FROM EMPLOYEE

WHERE DNO=5

SSN	LName	MName	FName	BirthDate	Address	Sex	Salary	SuperSSN	DNo
333445555	Nguyen	Thanh	Tung	12/08/1955	638 NVC Q5	Nam	40000	888665555	5
987987987	Nguyen	Manh	Hung	09/15/1962	Ba Ria VT	Nam	38000	333445555	5

 $\sigma_{\text{DNO}=5}$ (EMPLOYEE)

SELECT SSN, LNAME, MNAME, FNAME FROM EMPLOYEE WHERE DNO=5 AND SEX= 'Nam'

SSN	LNAME	MNAME	FNAME
333445555	Nguyen	Thanh	Tung
987987987	Nguyen	Manh	Hung

 $\pi_{\text{SSN, LNAME, MNAME, FNAME}}(\sigma_{\text{DNO}=5 \land \text{SEX}='\text{Nam'}}$ (EMPLOYEE))

Alias name

- SELECT SSN, LNAME AS 'Last Name', MNAME AS 'Middle Name', FNAME AS 'First Name'
- FROM EMPLOYEE

WHERE DNO=5 AND SEX= 'Nam'

SSN	LAST NAME	MIDDLE NAMI	FIRST NAME
333445555	Nguyen	Thanh	Tung
987987987	Nguyen	Manh	Hung

R1 \leftarrow ($\pi_{\text{SSN, LNAME, MNAME, FNAME}}(\sigma_{\text{DNO}=5 \land \text{SEX}='\text{Nam'}}(\text{EMPLOYEE})))$

 $\rho_{\text{result(SSN, LAST NAME, MIDDLE NAME, FIRST NAME}}(R1)$

Extension

SELECT SSN, LNAME + ' ' + MNAME+ ' ' + FNAME AS 'Full Name'

FROM EMPLOYEE

WHERE DNO=5 AND SEX= 'Nam'

SSN	Full Name
3334455555	Nguyen Thanh Tung
987987987	Nguyen Manh Hung

Extension

SELECT SSN, Salary*1.1 AS '10%SalaryIncrease'

FROM EMPLOYEE

WHERE DNo=5 AND Sex= 'Nam'

SSN	10%SalaryIncrease
333445555	33000
987987987	27500

 $\rho_{\text{SSN, 10\%SalaryIncrease}}(\pi_{\text{SSN, Salary*1.1}}(\sigma_{\text{DNo=5},\text{Sex='Nam'}}(\text{EMPLOYEE})))$

Duplicate tuples are eliminated

SELECT DIativiCT Salary FROM EMPLOYEE WHERE DNo=5 AND Sex= 'Nam'

Salary	
30000	
25000	- Cost
28000	Licore want to cap all tuples
38000	- Users want to see all tuples

Example

Find the SSN and first name of employees who work for the department 'Nghien cuu'

 $\begin{array}{l} \mathsf{R1} \leftarrow \mathsf{EMPLOYEE} \ \Join_{\mathsf{DNo=DNumber}} \ \mathsf{DEPARTMENT} \\ \mathsf{RESULT} \leftarrow \pi_{\mathsf{SSN, FName}}(\sigma_{\mathsf{DName='Nghien \, cuu'}}(\mathsf{R1})) \end{array}$

SELECTSSN, FNameFROMEMPLOYEE,DEPARTMENTDName= Nghien cuu'ANDDNo=DNumber



Priority

SELECT SSN, FName

FROM EMPLOYEE, DEPARTMENT

WHERE (DName= 'Nghien cuu' OR DName= 'Quan ly') AND DNo=DNumber

BETWEEN

SELECTSSN, FNameFROMEMPLOYEEWHERESalary>=20000 AND Salary=<30000</th>

SELECTSSN, FNameFROMEMPLOYEEWHERESalary BETWEEN 2000 AND 30000

NOT BETWEEN

SELECT SSN, FName

FROM EMPLOYEE

WHERE Salary NOT BETWEEN 20000 AND 30000
LIKE



NOT LIKE

SELECTSSN, FNameFROMEMPOYEEWHERELName LIKE 'Nguyen'

SELECTSSN, FNameFROMEMPLOYEEWHERELName NOT LIKE 'Nguyen'

ESCAPE

- SELECT SSN, FName
- FROM EMPOYEE
- WHERE Address LIKE '123x/%Nguyen' ESCAPE 'x'

Datetime

SELECT SSN, FName

FROM EMPLOYEE

WHERE BDate BETWEEN '1955-12-08' AND '1966-07-19'

'1955-12-08'YYYY-MM-DD'17:30:00'HH:MI:SS'12/08/1955'MM/DD/YYYY'05:30 PM''December 8, 1955'

'1955-12-08 17:30:00'

NULL

- SQL allows attributes to have value NULL
 - Value unknown
 - Value inapplicable
 - Value withheld
- Operation on a NULL and any value, the result is NULL
 - x has a value NULL
 - x + 3 is also NULL
 - x + 3 is not a legal SQL expression
- Comparison on a NULL value and any value, the result is UNKNOWN
 - The value of x = 3 is UNKNOWN
 - The comparison x = 3 is not correct SQL

NULL

SELECT SSN, FName

FROM EMPLOYEE

WHERE SuperSSN IS NULL

SELECTSSN, FNameFROMEMPLOYEE

WHERE SuperSSN IS NOT NULL

UNKNOWN

- The comparison involving NULL will result the threevalue logic
 - True (1)
 - False (0)
 - Unknown (1/2)
- Logical operator
 - x and y (Minimum value)
 - x or y (Maximum value)
 - not x (1-x)
- The condition in WHERE clauses will result *False* if tuples with *Unknown* value

FROM clause

Unspecified WHERE clause

SELECTSSN, DNumberFROMEMPLOYEE, DEPARTMENT

WHERE TRUE

SSN	DNumber
333445555	1
333445555	4
333445555	5
987987987	1
987987987	4
987987987	5

FROM clause

Alias name

SELECT DINtermee, Dilloocattion FROM DEDIARATIVIEWEINAS DEEP, TO EROTCATOOATSIONS AS DL WHERE DENDARDEDEDADIM Derumber

For each project locating in "Ha Noi", find its number, the department number that controls it, the last and first name of the manager, as well as his/her birth date and address

Find the last and first name of employees who work for the department 5 and work on the project "San pham X" with working hours are larger than 10

Find the first and last name of employees and their supervisor

Find the last and first name of employees who are supervised directly by "Nguyen Thanh Tung"

ORDER BY clause

- Is used for presenting a query in sorted order
- Syntax

SELECT <List_of_columns>

FROM <List_of_tables>

WHERE <Conditions>

ORDER BY<List_of columns>

- ASC (default)
- DESC

SELECTESSN, PNoFROMWORKS_ONORDERBY ESSN DESC, PNo

ESSN	PNo
999887777	10 🛉
999887777	30
987987987	10
987987987	30
987654321	10
987654321	20
987654321	30

Content

- Introduction
- Data definition
- Data manipulation
 - Basic queries
 - Set, set/multiset comparison and nested queries
 - Aggregate functions and grouping
 - Others
- Data update
- View definition

Index

Set operations in SQL

- SQL has implemented set operators
 - UNION
 - INTERSECT
 - EXCEPT
- The result is a set
 - Eliminate identical tuples
 - To keep identical tuples
 - UNION ALL
 - INTERSECT ALL
 - EXCEPT ALL

Set operations in SQL

Syntax

SELECT <Column_list> FROM <Table_list> WHERE <Condition> UNION [ALL]

SELECT <Column_list> FROM <Table_list> WHERE <Condition>

SELECT <Column_list> FROM <Table_list> WHERE <Condition> INTERSECT [ALL]

SELECT <Column_list> FROM <Table_list> WHERE <Condition>

SELECT <Column_list> FROM <Table_list> WHERE <Condition>
EXCEPT [ALL]

SELECT <Column_list> FROM <Table_list> WHERE <Condition>

- Find the project numbers that have
 - Either employees with the last name 'Nguyen',
 - Or been controlled by the department whose manager has the last name 'Nguyen'
 - SELECT PNo
 - FROM EMPLOYEE, WORKS_ON
 - WHERE SSN=ESSN AND LName= 'Nguyen'
 - UNION
 - SELECT PNumber
 - FROM EMPLOYEE, DEPARTMENT, PROJECT
 - WHERE SSN=MgrSSN AND DNumber=DNum AND LName= 'Nguyen'

Find employees who have dependents with the same name and sex

SELECT FName, Sex, SSN FROM EMPLOYEE

INTERSECT

SELECT Dependent_Name, Sex, ESSN FROM DEPENDENT

SELECT EM.*

FROM EMPLOYEE EM, DEPENDENT DE

WHERE EM.SSN=DE.ESSN

AND EM.FName=DE.Dependent_name

AND EM.Sex=DE.Sex

Example 6'

Find employees who have the same name and sex to dependents

SELECT FName, Sex FROM EMPLOYEE

INTERSECT

SELECT Dependent_Name, Sex FROM DEPENDENT

SELECT EM.* FROM EMPLOYEE EM, DEPENDENT DE WHERE EM.FName=DE.Dependent_Name AND EM.Sex=DE.Sex

Find employees who have no dependents

SELECT SSN FROM EMPLOYEE EXCEPT SELECT ESSN AS SSN FROM DEPENDENT

SELECT SSN, FName

FROM EMPLOYEE, DEPARTMENT

WHERE DName= 'Nghien cuu' AND DNo=DNumber

Outer query



- Queries can have several nested levels
 - Usually three
- Subqueries of a WHERE clause are connected by logical connective
 - OR, AND
- Subqueries will return
 - A single attribute and a single tuple (a single value)
 - A table (a set or multiset of tuples)

- WHERE clause of the outer query
 - <Expression> <set operation> <subquery>
 - Set comparison includes many operators
 - IN, NOT IN
 - ALL
 - ANY or SOME
 - Check whether the result of subqueries is empty or not
 - EXISTS
 - NOT EXISTS

- Categories
 - Subqueries that produce scalar values
 - WHERE clause of subqueries do not refer to attributes of relations in FROM clause of the outer query
 - Subqueries will be performed before the outer query, and be executed just one time
 - Correlated subqueries
 - WHERE clause of subqueries refer to at least one attribute of relations in FROM clause of the outer query
 - Subqueries will be executed many times, each time will correlate to one tuple of the outer query

Example – Scalar value subquery

SELECT SSN, FName

FROM EMPLOYEE, DEPT_LOCATIONS

WHERE DLocation= 'TP HCM' AND DNo=DNumber



SELECT WO.PNo

FROM EMLOYEE EM, WORKS_ON WO

WHERE EM.SSN=WO.ESSN AND EM.LName= 'Nguyen'

UNION

SELECT PR.PNumber

FROM EMPLOYEE EM, DEPARTMENT DE, PROJECT PR

WHERE EM.SSN=DE.MgrSSN AND DE.DNumber=PR.DNum

AND EM.LName= 'Nguyen'



Find employees who have no dependents

SELECT * FROM EMPLOYEE WHERE SSN NOT IN (SELECT ESSN **DEPENDENT**) FROM SELECT * FROM EMPLOYEE WHERE SSN <> ALL (SELECT ESSN FROM DEPENDENT)

Find employees whose salary is greater than <u>at</u> <u>least one</u> salary of employees in department 4

```
SELECT *
FROM EMPLOYEE
WHERE Salary > ANY(
    Salary > ANY(
    SELECT Salary
    FROM EMPLOYEE
    WHERE DNo=4)
SELECT EM1.*
```

- FROM EMPLOYEE EM1, EMPLOYEE EM2
- WHERE EM1.Salary > EM2.Salary AND EM2.DNo=4

Find employees whose salary is greater than all salaries of employees in the department 4

SELECT*FROMEMPLOYEEWHERESalary > ALL (SELECTSalaryFROMSalaryFROMEMPLOYEEWHEREDNo=4)

Find managers who have at least one dependent

SELECT *FROMEMPLOYEEWHERESSN IN (SELECT ESSN FROM DEPENDENT)AND SSN IN (SELECT MgrSSN FROM DEPARTMENT)

Example - Correlated subquery

SELECT SSN, FName

FROM EMPLOYEE, DEPARTMENT

WHERE DName= 'Nghien cuu' AND DNo=DNumber



Retrieve employees who have dependents with the same first name and same sex as the employees

```
SELECT *
FROM EMPLOYEE EM
WHERE EXISTS (
SELECT *
FROM DEPENDENT DE
WHERE EM.SSN=DE.SSN
AND EM.FName=DE.Name
AND EM.Sex=DE.Sex )
```

Find employees who have no dependents

SELECT *
FROM EMPLOYEE
WHERE NOT EXISTS (
SELECT *
FROM DEPENDENT
WHERE SSN=ESSN)
Find employees whose salary is greater than <u>at</u> <u>least one</u> salary of employees in department 4

SELECT *
FROM EMPLOYEE EM1
WHERE EXISTS (
 SELECT *
 FROM EMPLOYEE EM2
 WHERE EM2.PNo=4
 AND EM1.Salary>EM2.Salary)

Find managers who have <u>at least one</u> dependent

SELECT * FROM EMPLOYEE WHERE EXISTS (**SELECT** * FROM DEPENDENT WHERE SSN=ESSN) AND EXISTS (SELECT * FROM DEPARTMENT WHERE SSN=MgrSSN)

Discussion IN and EXISTS

IN

- <Column_name> IN <Subquery>
- Attributes in the subquery's SELECT clause have the same data types as attributes in the outer query's WHERE clause
- EXISTS
 - Do not need attributes, constants or any expressions before it
 - Do not need to specify column names in the subquery's SELECT clause
 - Queries containing "= ANY" or IN can be converted queries containing EXISTS

Discussion

- Comparison of one value and members in a set
 - any/some or exists of nested queries ⇔ equijoin of basic queries

Divide operation in SQL





R÷S	А	В	С
a,	α	а	γ
·	γ	а	γ

R÷S is a set of values a_i in R such that there is <u>no</u> values b_i in S that makes the tuple (a_i, b_i) does <u>not</u> exist in R

- Retrieve the first name of employees who work on all projects
 - Retrieve the first name of employees such that there is no projects that they do not work on
 - R: WORKS_ON(ESSN, PNo)
 - S: PROJECT(PNumber)
 - R÷S: RESULT(ESSN)
 - Joining RESULT to EMPLOYEE to retrieve FName

Using NOT EXISTS two times

SELECT R1.A, R1.B, R1.C FROM **R** R1 WHERE NOT EXISTS (SELECT * FROM S WHERE NOT EXISTS (SELECT * FROM R R2 WHERE R2.D=S.D AND R2.E=S.E AND R1.A=R2.A AND R1.B=R2.B AND R1.C=R2.C))

SELECT EM.FName

FROM EMPLOYEE EM, WORKS_ON WO1

WHERE EM.SSN=WO1.ESSN

AND NOT EXISTS (

SELECT *

FROM PROJECT PR

WHERE NOT EXISTS (

SELECT *

FROM WORKS_ON WO2

WHERE WO2.PNo=PR.PNumber

AND WO1.ESSN=WO2.ESSN))

Using NOT EXISTS and EXCEPT

SELECT R1.A, R1.B, R1.C

FROM R R1

WHERE NOT EXISTS (

(SELECT D, E FROM S)

EXCEPT

(SELECT D, E FROM R R2 WHERE R1.A=R2.A

AND R1.B=R2.B

AND R1.C=R2.C))

SELECT FName

FROM EMPLOYEE

WHERE NOT EXISTS (

(SELECT PNumber FROM PROJECT)

EXCEPT

(SELECT PNo FROM WORKS_ON WHERE SSN=ESSN))

Content

- Introduction
- Data definition
- Data manipulation
 - Basic queries
 - Set, set/multiset comparison and nested queries
 - Aggregate functions and grouping
 - Others
- Data update
- View definition

Index

Aggregate functions

- COUNT
 - COUNT(*) : the number of rows
 - COUNT(<Column_name>): the number of non-zero values of the column
 - COUNT(DISTINCT <Column_name>): the number of different and non-zero values of the column
- MIN
- MAX
- SUM
- AVG

These function is in SELECT clause

Find the sum of salary, highest salary, lowest salary and average salary of employees

SELECT SUM(Salary), MAX(Salary), MIN(Salary), AVG(Salary)FROM EMPLOYEE

Find the number of employees in the department 'Nghien cuu'

SELECT COUNT(*) AS Num_Emp

FROM EMPLOYEE, DEPARTMENT

WHERE PNo=PNumber AND PName= 'Nghien cuu'

Find the number of employees for each department

DNo	Num_Emp
5	3
4	3
1	1

SSN	LName	MName	FName	BirthDate	Address	Sex	Salary	SuperSSN	DNo
333445555	Nguyen	Thanh	Tung	12/08/1955	638 NVC Q5	Nam	40000	888665555	5
987987987	Nguyen	Manh	Hung	09/15/1962	Ba Ria VT	Nam	38000	333445555	5
453453453	Tran	Thanh	Tam	07/31/1972	543 MTL Q1	Nu	25000	333445555	5
999887777	Bui	Ngoc	Hang	07/19/1968	33 NTH Q1	Nu	38000	987654321	4
987654321	Le	Quynh	Nhu	07620/1951	219 TD Q3	Nu	43000	888665555	4
987987987	Tran	Hong	Quang	04/08/1969	980 LHP Q5	Nam	25000	987654321	4
888665555	Pham	Van	Vinh	11/10/1945	450 TV HN	Nam	55000	NULL	1

Grouping

Syntax

SELECT <List_of_columns>

FROM <List_of_ tables>

WHERE <Conditions>

GROUP BY <List_of_grouping_columns>

- After grouping
 - Each group will have identical values at grouping attributes

Find the number of employees for each department

SELECTDNo, COUNT(*) AS Num_EmpFROMEMPLOYEEGROUP BY DNo

SELECT DName, COUNT(*) AS Num_EmpFROM EMPLOYEE, DEPARTMENTWHERE DNo=DNumberGROUP BY DName

For each employee, retrieve the SSN, first name, last name, number of projects as well as total of hours that the employee works on

ESSN	PNo	Hours
123456789	1	32.5
123456789	2	7.5
333445555	2	10.0
333445555	3	10.0
333445555	10	10.0
888665555	20	20.0
987987987	10	35.0
987987987	30	5.0
987654321	30	20.0
987654321	20	15.0
453453453	1	20.0
453453453	2	20.0

SELECT ESSN, COUNT(*) AS Num_Pro, SUM(Hours) AS Total_Hours FROM WORKS_ON GROUP BY ESSN

SELECT LName, FName, COUNT(*) AS Num_Pro,

SUM(THOIGIAN) AS Total_Hours

FROM WORKS_ON, EMPLOYEE

WHERE ESSN=SSN

GROUP BY ESSN, LName, FName

Find employees who work on two or more projects

ESSN	DNo	Hours	
123456789	1	32.5	
123456789	2	7.5	
333445555	2	10.0	
333445555	3	10.0	
3334455555	10	10.0	
888665555	20	20.0	Eliminated
987987987	10	35.0	
987987987	30	5.0	
987654321	30	20.0	
987654321	20	15.0	
453453453	1	20.0	
453453453	2	20.0	

Conditions on groups

Syntax

SELECT <List_of_columns>

FROM <List_of_tables>

WHERE <Conditions>

GROUP BY <List_of_grouping_columns>

HAVING <Conditions>

Find employees who work on two or more projects

SELECT ESSN FROM WORKS_ON GROUP BY ESSN HAVING COUNT(*) >= 2

Retrieve the department name whose the average salary of employees that is greater than 20000

SELECT DNo, AVG(Salary) AS Avg_Salary

FROM EMPLOYEE

GROUP BY DNo

HAVING AVG(Salary) > 20000

SELECT DName, AVG(Salary) AS Avg_Salary

FROM EMPLOYEE, DEPARTMENT

WHERE DNo=DNumber

GROUP BY DNo, DName

HAVING AVG(Salary) > 20000

Discussion

- GROUP BY
 - Attributes in SELECT clause (excepting attributes of aggregate functions) must appear in GROUP BY clause

HAVING

- Use aggregate functions in SELECT clause to check a certain condition
- Just validate the conditions for groups, not a condition for filtering rows
- After grouping, conditions on groups will be performed

Discussion

- The order of the query execution
 - (1) Pick out rows that satify conditions in WHERE clause
 - (2) Group these rows into many groups in GROUP BY clause
 - (3) Apply aggregate functions for each group
 - (4) Eliminate groups that do not satisfy conditions in HAVING clause
 - (5) Retrieve values from columns and aggregate functions in SELECT clause

Find departments that have the highest average salary

SELECT DNo, AVG(Salary) AS Avg_Salary FROM EMPLOYEE GROUP BY DNo HAVING MXX (Sala(Salary))LL (SELECT AVG(Salary) FROM EMPLOYEE GROUP BY DNo)

Find three employees who have the highest salary

SELECT FName FROM EMPLOYEE EM1 WHERE 2 >= (SELECT COUNT(*) FROM EMPLOYEE EM2 WHERE EM2.Salary>EM1.Salary)

Discussion

- Find three employees who have the highest salary
 - If salaries are redundant, then ???

Find the first name of employees who work on <u>all</u> projects

SELECT SSN, FName FROM EMPLOYEE, WORKS_ON WHERE SSN=ESSN GROUP BY SSN, FName HAVING COUNT(*) = (SELECT COUNT(*) FROM PROJECT)

Content

- Introduction
- Data definition
- Data manipulation
 - Basic queries
 - Set, set/multiset comparison and nested queries
 - Aggregate functions and grouping
 - Other queries
- Data update
- View definition

Index

Other queries

- Subquery in FROM clause
- Joining conditions in FROM clause
 - Natural join
 - Outer join
- CASE structure

Subquery in FROM clause

- The result of a subquery is a table
 - Intermediate table in the process of query execution
 - Do not store this result into the database

Syntax

SELECT <List_of_columns>

FROM R1, R2, (<Subquery>) **AS** Table_name

WHERE <Conditions>

Retrieve the department name whose the average salary of employees that is greater than 20000

SELECT DNo, AVG(Salary) AS Avg_Salary

FROM EMPLOYEE

GROUP BY DNo

HAVING AVG(Salary) > 20000

SELECT DName, AVG(Salary) AS Avg_Salary

FROM EMPLOYEE, DEPARTMENT

WHERE DNo=DNumber

GROUP BY DNo, DName

HAVING AVG(Salary) > 20000

Retrieve the department name whose the average salary of employees that is greater than 20000

SELECT DName, TEMP.Avg_Salary FROM DEPARTMENT, (SELECT DNo, AVG(Salary) AS Avg_Salary FROM EMPLOYEE GROUP BY DNo HAVING AVG(Salary)> 20000) AS TEMP

WHERE DNumber=TEMP.DNo

Join conditions in FROM clause

Equijoin

SELECT <List_of_columns>

FROM R1 [INNER] JOIN R2 ON <Expression>

WHERE <Conditions>

Outer join

SELECT <List_of_columns>

FROM R1 LEFT | RIGHT [OUTER] JOIN R2 ON <Expression>

WHERE <Conditions>
Retrieve the SSN and first name of employees who work for the department 'Nghien cuu'

SELECT SSN, FName

- FROM EMPLOYEE, DEPARTMENT
- WHERE DName= 'Nghien cuu' AND DNo=DNumber
- SELECT SSN, FName
- FROM EMPLOYEE INNER JOIN DEPARTMENT ON DNo=DNumber
- WHERE DName= 'Nghien cuu'

Retrieve the fist and last name of employees, as well as the department name that they are managers

FName	LName	DName
Tung	Nguyen	Nghien cuu
Hang	Bui	null
Nhu	Le	null
Vinh	Pham	Quan ly

SELECTFName, LName, DNameFROMEMPLOYEE, DEPARTMENTWHERESSN=MgrSSN



FROM EMPLOYEE LEFT JOIN DEPARTMENT ON SSN=MgrSSN



FROM DEPARTMENT RIGHT JOIN EMPLOYEE ON SSN=MgrSSN

Retrieve the first and last name of employees, the name of projects that they work on (<u>if any</u>)



SELECT EM.FName, PR.PName

FROM (WORKS_ON WO JOIN PROJECT PR ON PNO=PNumber) RIGHT JOIN EMPLOYEE EM ON WO.ESSN=EM.SSN

CASE structure

Allow us to check conditions or output the information in each case

Syntax

```
CASE <Column_name>
WHEN <Value> THEN <Expression>
WHEN <Value> THEN <Expression>
...
[ELSE <Expression>]
END
```

Retrieve the last and first name of employees who are gonna reach the retired age (male: 60 years old, female: 55 years old)

```
SELECT LName, FName
FROM EMPLOYEE
WHERE YEAR(GETDATE()) - YEAR(BirthDate) >= ( CASE Sex
WHEN 'Nam' THEN 60
WHEN 'Nu' THEN 55
END )
```

Retrieve the last and first name of employees, as well as their retired year

SELECT LName, FName,

(CASE Sex

WHEN 'Nam' THEN YEAR(BirthDate) + 60

WHEN 'Nu' THEN YEAR(BirthDate) + 55

END) AS RetiredYear

FROM EMPLOYEE

Summary

SELECT <List_of_columns>

FROM <List_of_tables>

[WHERE <Conditions>]

[**GROUP BY** <List_grouping_columns>]

[HAVING <Conditions>]

[**ORDER BY** <List_of_ordering_columns>]

Content

- Introduction
- Data definition
- Data manipulation
- Data update
 - Insert
 - Delete
 - Update
- View definition
- Index

INSERT command

- Is used to add 1 or more tuple(s) to a relation
- In order to add a tuple
 - Relation name
 - List of column names
 - List of values for the tuple

INSERT command

Syntax (one tuple)

INSERT INTO <Table_name>(<List_of_columns>)

VALUES (<List_of_values>)

```
INSERT INTO EMPLOYEE(LName, MName, FName, SSN)
VALUES ('Le', 'Van', 'Tuyen', '635635635')
```

```
INSERT INTO EMPLOYEE(LName, MName, FName, SSN, Address)
VALUES ('Le', 'Van', 'Tuyen', '635635635', NULL)
```

```
INSERT INTO EMPLOYEE
VALUES ('Le', 'Van', 'Tuyen', '635635635', '12/30/1952', '98 HV', 'Nam',
'37000', 4)
```

INSERT command

- Discussion
 - The order of values is the same to the order of columns
 - The NULL value can be used for non-primary-key attributes
 - INSERT command will raise errors if the integrity constraint is violated
 - Primary key
 - Reference
 - NOT NULL constraint

INSERT command

Syntax (many tuples)

INSERT INTO <Table_name>(<List_of_columns>)

<Query>

CREATE TABLE ANALYSE_DEPT (DName VARCHAR(20), Number_Emp INT, Total_Salary INT

INSERT INTO ANALYSE_DEPT(DName, Number_Emp, Total_Salary) SELECT DName, COUNT(SSN), SUM(Salary) FROM EMPLOYEE, DEPARTMENT WHERE DNo=DNumber GROUP BY DName

Is used to remove tuples from a relation



DELETE FROM <Table_name>

[WHERE <Conditions>]

DELETE FROM EMPLOYEE WHERE LName= 'Tran'

DELETE FROM EMPLOYEE WHERE SSN= '345345345'

DELETE FROM EMPLOYEE

Remove employees who work for the department 'Nghien cuu'

> DELETE FROM EMPLOYEE WHERE DNo IN (SELECT DNumber FROM DEPARTMENT WHERE DName= 'Nghien cuu')

- Discussion
 - The number of removed tuples depends on the condition in WHERE clause
 - A missing WHERE clause specifies that all tuples can be deleted
 - DELETE command can cause the violation of reference constraints
 - Do not permit to remove
 - Remove tuples whose value is being referred
 - * CASCADE
 - Set the NULL value to cho những giá trị tham chiếu

SSN	LName	MName	FName	BirthDate	Address	Sex	Salary	SuperSSN	DNo
3334455555	Nguyen	Thanh	Tung	12/08/1955	638 NVC Q5	Nam	40000	888665555	5
987987987	Nguyen	Manh	Hung	09/15/1962	Ba Ria VT	Nam	38000	333445555	5
453453453	Tran	Thanh	Tam	07/31/1972	543 MTL Q1	Nu	25000	333445555	5
999887777	Bui	Ngoc	Hang	07/19/1968	33 NTH Q1	Nu	38000	987654321	4
987654321	Le	Quynh	Nhu	07620/1951	219 TD Q3	Nu	43000	888665555	4
987987987	Tran	Hong	Quang	04/08/1969	980 LHP Q5	Nam	25000	987654321	4
888665555	Pham	Van	Vinh	11/10/1945	450 TV HN	Nam	55000	NULL	1

ESSN	DNo	Hours
333445555	10	10.0
888665555	20	20.0
987987987	10	35.0
987987987	30	5.0
987654321	30	20.0
453453453	1	20.0

DName	DNumber	MgrSSN	MgrStartDate
Nghien cuu	5	333445555	05/22/1988
Dieu hanh	4	987987987	01/01/1995
Quan ly	1	888665555	06/19/1981

SSN	LName	MName	FName	BirthDate	Address	Sex	Salary	SuperSSN	DNo
333445555	Nguyen	Thanh	Tung	12/08/1955	638 NVC Q5	Nam	40000	888665555	NUSLL
987987987	Nguyen	Manh	Hung	09/15/1962	Ba Ria VT	Nam	38000	333445555	NUSLL
453453453	Tran	Thanh	Tam	07/31/1972	543 MTL Q1	Nu	25000	333445555	NUSLL
999887777	Bui	Ngoc	Hang	07/19/1968	33 NTH Q1	Nu	38000	987654321	4
987654321	Le	Quynh	Nhu	07620/1951	219 TD Q3	Nu	43000	888665555	4
987987987	Tran	Hong	Quang	04/08/1969	980 LHP Q5	Nam	25000	987654321	4
888665555	Pham	Van	Vinh	11/10/1945	450 TV HN	Nam	55000	NULL	1

UPDATE command

Is used to change the value of attributes

Syntax

UPDATE <Table_name>

SET <Attribute_name>=<The_new_value>,

<Attribute_name>=<The_new_value>,

•••

[WHERE <Condition>]

UPDATE EMPLOYEESET BDate='08/12/1965'WHERE SSN='333445555'

UPDATE EMPLOYEE

SET Salary=Salary*1.1

Change the location and controlling department number of the project 10 to 'Vung Tau' and 5, respectively

UPDATE PROJECTSET PLocation='Vung Tau', DNum=5WHERE PNumber=10

UPDATE command

- Discussion
 - Tuples that satisfy conditions in WHERE clause will be modified to the new value
 - A missing WHERE clause specifies that all tuples can be modified
 - UPDATE command can cause violations of the reference constraint
 - Do not allow to modify
 - Modify the values of tuples that are being referred
 - * CASCADE

Content

- Introduction
- Data definition
- Data manipulation
- Data update
- View definition
 - Definition
 - Query
 - Modification
- Index

View

- Table is a relation that exist actually the database
 - Stored in some physical organization
 - Persistent
- View is also a relation
 - Do not exist physically (virtual table)
 - Do not contain the data
 - Is derived from other tables
 - Can query or even modify the data through views

View

- Why do we use views?
 - Hide the complexity of data
 - Simplify queries
 - Present data with the purpose "easy to use"
 - Mechanism of data safety

Definition



CREATE VIEW <View_name> AS

<Query>

DROP VIEW <View_name>

- View contains
 - A list of attributes that are the same as attributes in SELECT clause
 - The number of tuples depending on the conditions in WHERE clause
 - Data derived from tables in FROM clause

CREATE VIEW EMP_DEP5 AS

SELECT SSN, LName, MName, FName

FROM EMPLOYEE

WHERE DNo=5

CREATE VIEW STATISTIC_DEP AS

SELECT DNo, DName, COUNT(*) AS Number_Emp, SUM(Salary) AS Total_Salary

FROM EMPLOYEE, DEPARTMENT

WHERE DNo=DNumber

GROUP BY DNumber, DName

Querying views

Although views do not contain data, we can do the query on views

SELECT FName

FROM EMP_DEP5

WHERE LName LIKE 'Nguyen'

 $\mathsf{EMP_DEP5} \leftarrow \pi_{\mathsf{SSN, LName, MName, FName}}(\sigma_{\mathsf{DNo=5}}(\mathsf{EMPLOYEE}))$

$$\pi_{\text{FName}}(\sigma_{\text{LName='Nguyen'}} \text{ (EMP_DEP5))}$$

Querying views

Can query data from both tables and views

SELECT LName, FName, PName, Hours FROM EMP_DEP5, WORKS_ON, PROJECT WHERE SSN=ESSN AND PNo=PNumber

 $\mathsf{EMP_DEP5} \leftarrow \pi_{\mathsf{SSN, LName, MName, FName}}(\sigma_{\mathsf{DNo=5}}(\mathsf{EMPLOYEE}))$

 $\mathsf{TMP} \leftarrow \mathsf{EMP_DEP5} \bowtie_{\mathsf{SSN}=\mathsf{ESSN}}(\mathsf{WORKS_ON} \bowtie_{\mathsf{PNo}=\mathsf{PNumber}} \mathsf{PROJECT})$ $\pi_{\mathsf{LName},\mathsf{FName},\mathsf{PName},\mathsf{Hours}}(\mathsf{TMP})$

Modifying views

- Can apply INSERT, DELETE, and UPDATE commands to <u>simple views</u>
 - Views built on one table and having the key attribute of that table
- Cannot modify views
 - Views have a key word DISTINCT
 - Views use aggregate functions
 - Views have extended SELECT clause
 - Views are derived from table containing constrains on columns
 - Views are derived from many tables

Modifying views

Modify the last name of employee '123456789' in department 5 to 'Pham'

UPDATE EMP_DEP5 SET LName= 'Pham' WHERE SSN= '123456789'

Content

- Introduction
- Data definition
- Data manipulation
- Data update
- View definition

Index
Index

The index on an attribute A is the data structure that makes it efficient to find tuples having a fixed value for attribute A

SELECT * FROM EMPLOYEE WHERE DNo=5 AND Sex= 'Nu'

Read 10.000 tuples

Read 200 tuples

Table EMPLOYEE has 10.000 tuples

There are 200 employees who work for the department 5

Read 70 tuples

Index



CREATE INDEX <Index_name> **ON** <Table_name>(<Column_name>)

DROP INDEX <Index_name>



CREATE INDEX DNo_IND ON EMPLOYEE(DNo)

CREATE INDEX DNo_Sex_IND ON EMPLOYEE(DNo, Sex)

Index

- Discussion
 - Speed up queries in which a value for an attribute is specified, join operations
 - Make insertion, deletion, and update more complex and time-consuming
 - Cost
 - Index storage
 - Disk access (HDD)
- Selection of indexes
 - One of the principal factors that influence a database
 - One of the hardest parts of database design

Example

- Examine the relation WORKS_ON(ESSN, PNo, Hours)
- Assume that
- WORKS_ON is stored in 10 disk blocks
 - The cost of examining the entire relation WORKS_ON is 10
- On the average, an employee works on 3 projects and a project has 3 employees
 - Tuples for a given employee or project are likely to be spread over the 10 disk blocks
 - The cost of finding 3 tuples for an employee or a project will take 3 disk accesses (if we have indexes)
- Using indexes to locate tuples
 - The cost of reading an index's block is 1 disk access
- Insertion costs 2 disk accesses

Example

- Suppose that the following queries are performed frequently
 - Q1 SELECT PNo, Hours

FROM WORKS_ON

WHERE ESSN= '123456789'

- Q2 SELECT ESSN FROM WORKS_ON WHERE PNo=1 AND Hours=20.5
- Q3 INSERT INTO WORKS_ON VALUES (123456789' , 1, 20.5)

Example

The costs of each of the 3 queries

Queries	No index	ESSN index	PNo index	Both indexes
Q1	10	4	10	4
Q2	10	10	4	4
Q3	2	4	4	6
Average cost	2 + 8p1 + 8p2	4 + 6p2	4 + 6p1	6 - 2p1 – 2p2

The fraction of the time we do Q1 is p1 The fraction of the time we do Q2 is p2 The fraction of the time we do Q3 là 1 - p1 - p2

