

ARP and RARP

Objectives

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Upon completion you will be able to:

- *Understand the need for ARP*
- *Understand the cases in which ARP is used*
- *Understand the components and interactions in an ARP package*
- *Understand the need for RARP*

Figure 7.1 *ARP and RARP*

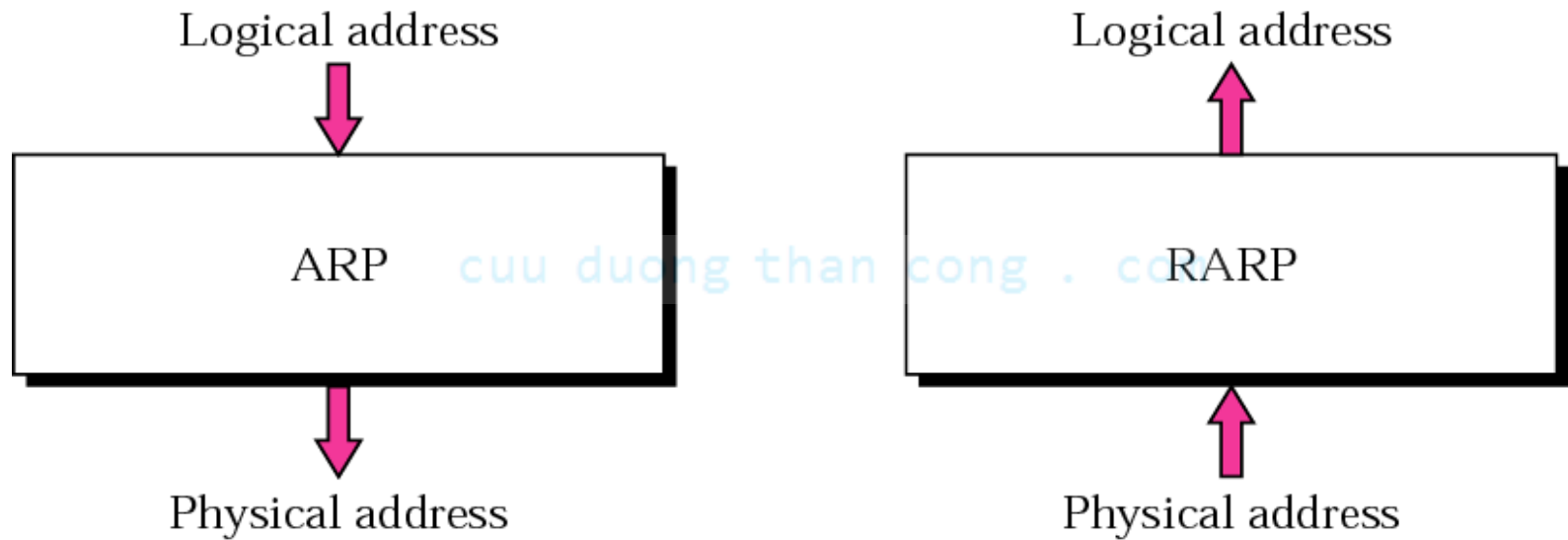
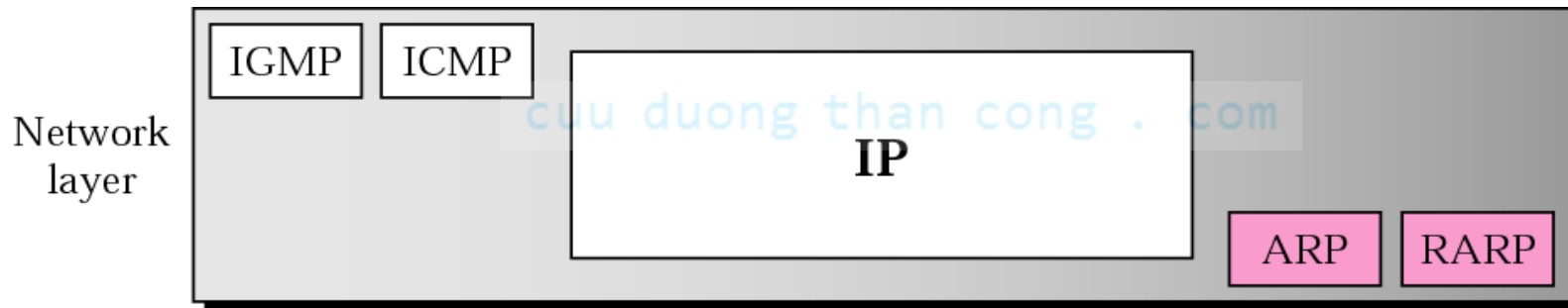


Figure 7.2 *Position of ARP and RARP in TCP/IP protocol suite*



7.1 ARP

ARP associates an IP address with its physical address. On a typical physical network, such as a LAN, each device on a link is identified by a physical or station address that is usually imprinted on the NIC.

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The topics discussed in this section include:

Packet Format

Encapsulation

Operation

ARP over ATM

Proxy ARP

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Figure 7.3 *ARP operation*

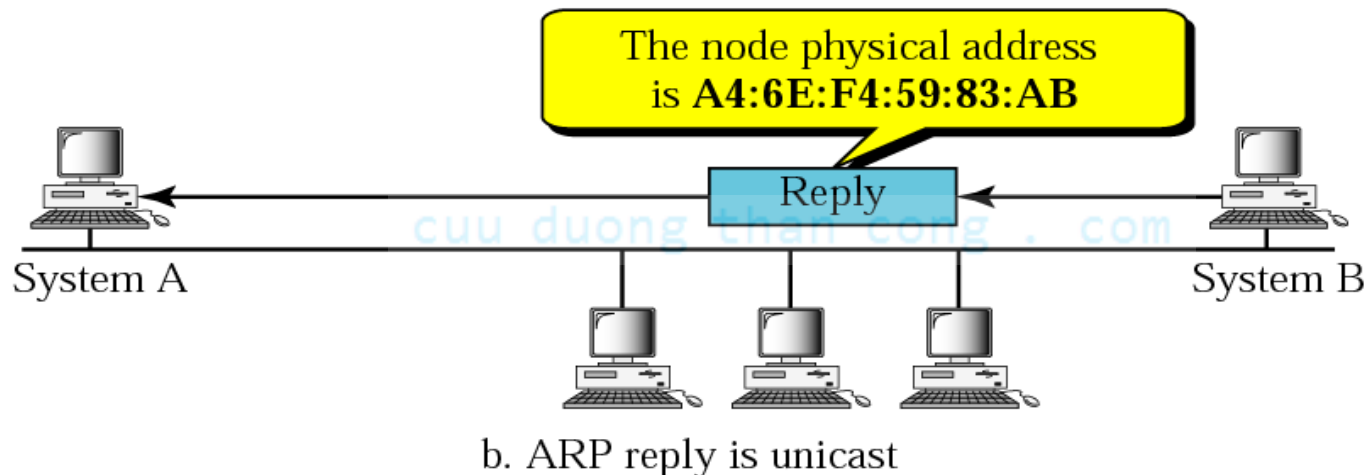
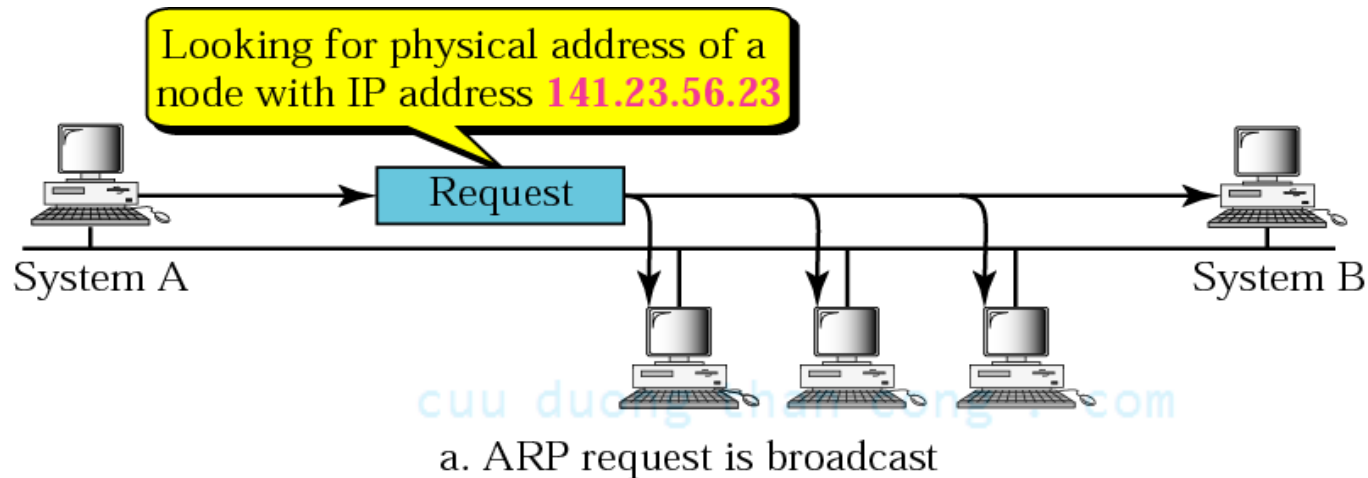


Figure 7.4 *ARP packet*

Hardware Type		Protocol Type
Hardware length	Protocol length	Operation Request 1, Reply 2
Sender hardware address (For example, 6 bytes for Ethernet)		
Sender protocol address (For example, 4 bytes for IP)		
Target hardware address (For example, 6 bytes for Ethernet) (It is not filled in a request)		
Target protocol address (For example, 4 bytes for IP)		

Figure 7.5 *Encapsulation of ARP packet*

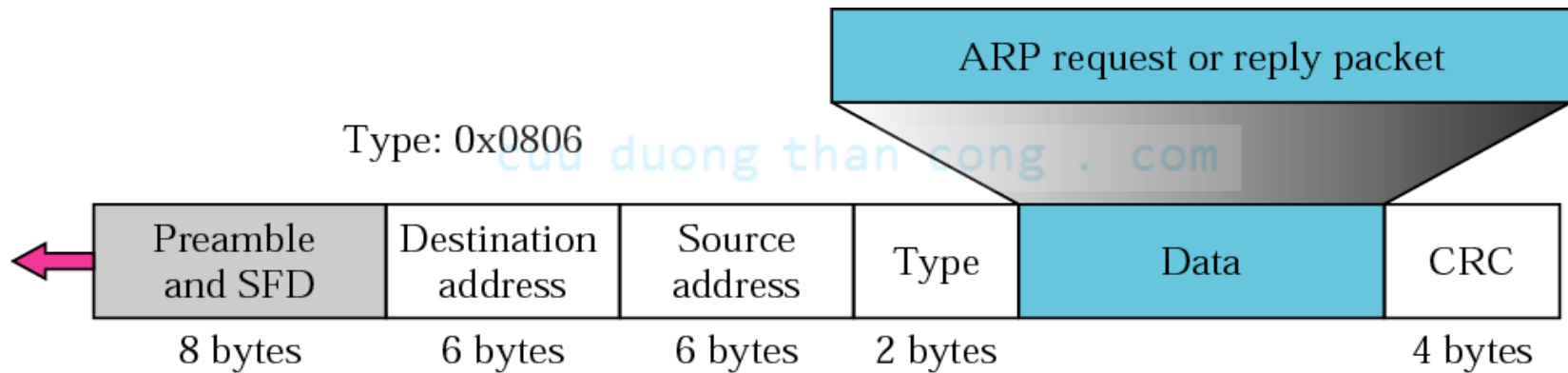
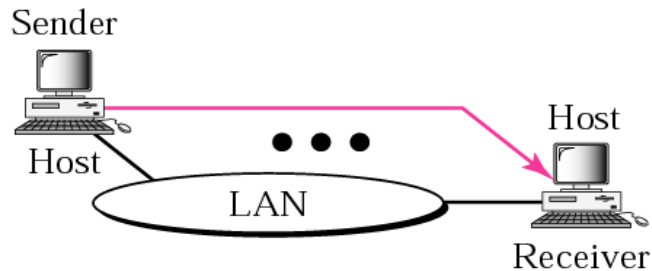


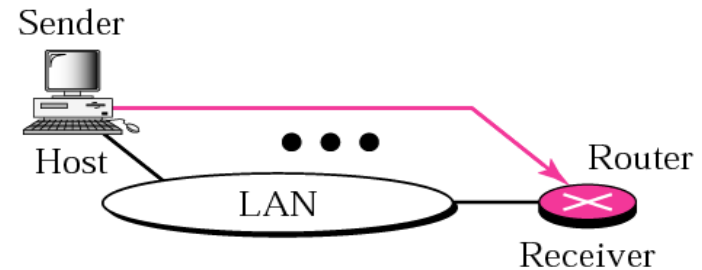
Figure 7.6 *Four cases using ARP*

Target IP address:
Destination address in the IP datagram



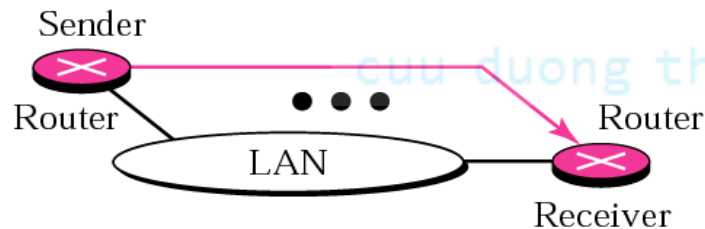
Case 1. A host has a packet to send to another host on the same network.

Target IP address:
IP address of a router



Case 2. A host wants to send a packet to another host on another network. It must first be delivered to a router.

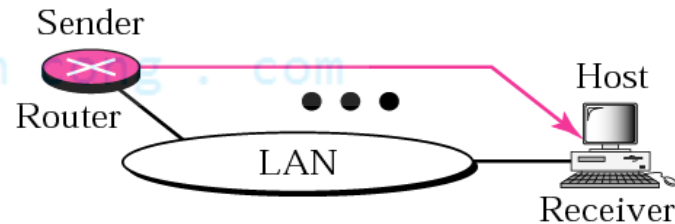
Target IP address:
IP address of the appropriate router
found in the routing table



Case 3. A router receives a packet to be sent to a host on another network.

It must first be delivered to the appropriate router.

Target IP address:
Destination address in the IP datagram



Case 4. A router receives a packet to be sent to a host on the same network.



Note:

*An ARP request is broadcast;
an ARP reply is unicast.*

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EXAMPLE 1

*A host with IP address **130.23.43.20** and physical address **B2:34:55:10:22:10** has a packet to send to another host with IP address **130.23.43.25** and physical address **A4:6E:F4:59:83:AB** (which is unknown to the first host). The two hosts are on the same Ethernet network. Show the ARP request and reply packets encapsulated in Ethernet frames.*

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EXAMPLE 1 (CONTINUED)

Solution

Figure 7.7 shows the ARP request and reply packets. Note that the ARP data field in this case is 28 bytes, and that the individual addresses do not fit in the 4-byte boundary. That is why we do not show the regular 4-byte boundaries for these addresses. Also note that the IP addresses are shown in hexadecimal. For information on binary or hexadecimal notation see Appendix B.

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Figure 7.7 Example 1

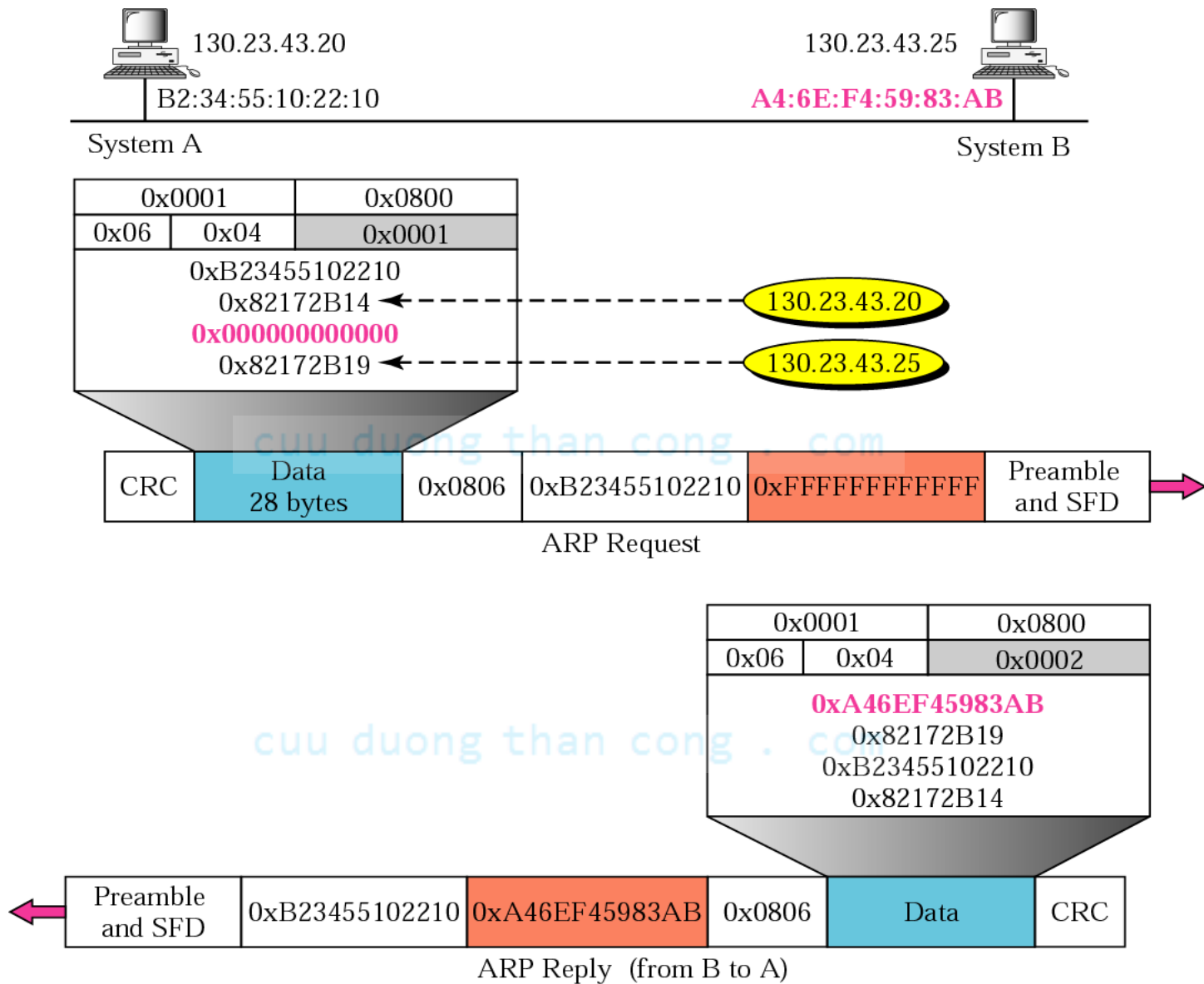
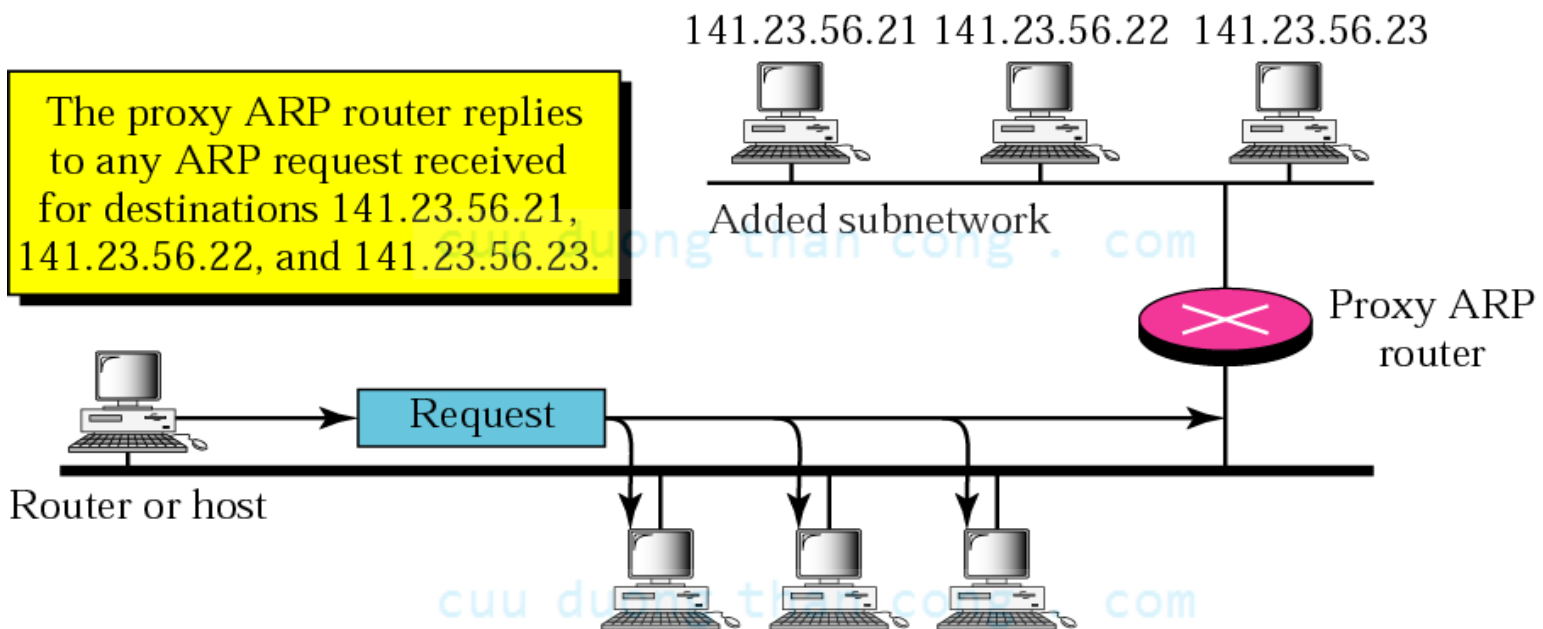


Figure 7.8 *Proxy ARP*



7.2 ARP PACKAGE

In this section, we give an example of a simplified ARP software package to show the components and the relationships between the components. This ARP package involves five modules: a cache table, queues, an output module, an input module, and a cache-control module.

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The topics discussed in this section include:

Cache Table

Queues

Output Module

Input Module

Cache-Control Module

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Figure 7.9 *ARP components*

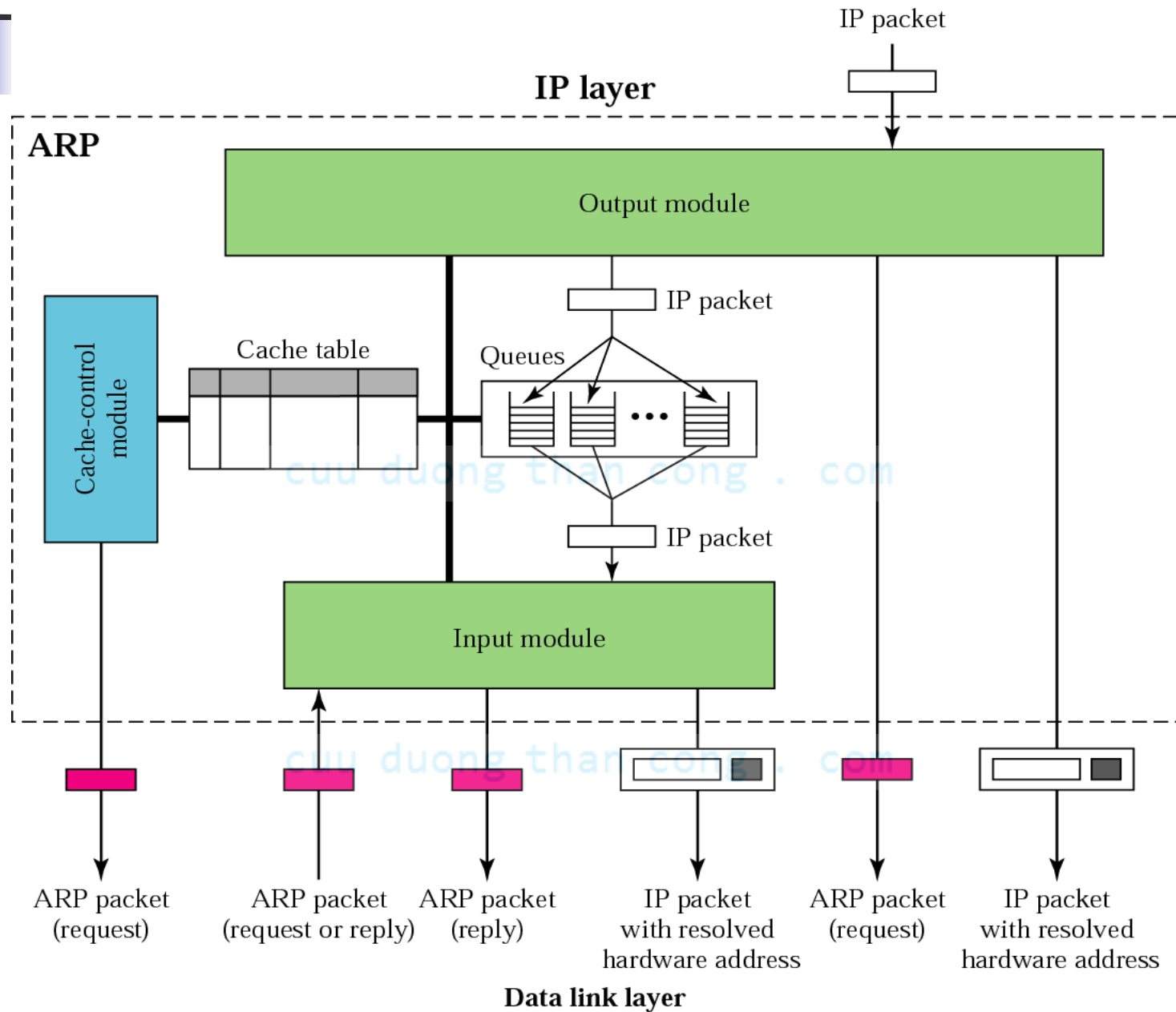


Table 7.1 *Original cache table used for examples*

<i>State</i>	<i>Queue</i>	<i>Attempt</i>	<i>Time-Out</i>	<i>Protocol Addr.</i>	<i>Hardware Addr.</i>
R	5		900	180.3.6.1	ACAE32457342
P	2	2		129.34.4.8	
P	14	5		201.11.56.7	
R	8		450	114.5.7.89	457342ACAE32
P	12	1		220.55.5.7	
F					
R	9		60	19.1.7.82	4573E3242ACA
P	18	3		188.11.8.71	



EXAMPLE 2

The ARP output module receives an IP datagram (from the IP layer) with the destination address 114.5.7.89. It checks the cache table and finds that an entry exists for this destination with the RESOLVED state (R in the table). It extracts the hardware address, which is 457342ACAE32, and sends the packet and the address to the data link layer for transmission. The cache table remains the same.



EXAMPLE 3

Twenty seconds later, the ARP output module receives an IP datagram (from the IP layer) with the destination address 116.1.7.22. It checks the cache table and does not find this destination in the table. The module adds an entry to the table with the state PENDING and the Attempt value 1. It creates a new queue for this destination and enqueues the packet. It then sends an ARP request to the data link layer for this destination. The new cache table is shown in Table 7.2.

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Table 7.2 Updated cache table for Example 3

<i>State</i>	<i>Queue</i>	<i>Attempt</i>	<i>Time-Out</i>	<i>Protocol Addr.</i>	<i>Hardware Addr.</i>
R	5		900	180.3.6.1	ACAE32457342
P	2	2		129.34.4.8	
P	14	5		201.11.56.7	
R	8		450	114.5.7.89	457342ACAE32
P	12	1		220.55.5.7	
P	23	1		116.1.7.22	
R	9	60		19.1.7.82	4573E3242ACA
P	18	3		188.11.8.71	



EXAMPLE 4

*Fifteen seconds later, the ARP input module receives an ARP packet with target protocol (IP) address 188.11.8.71. The module checks the table and finds this address. It changes the state of the entry to **RESOLVED** and sets the time-out value to 900. The module then adds the target hardware address (E34573242ACA) to the entry. Now it accesses queue 18 and sends all the packets in this queue, one by one, to the data link layer. The new cache table is shown in Table 7.3.*

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Table 7.3 Updated cache table for Example 4

<i>State</i>	<i>Queue</i>	<i>Attempt</i>	<i>Time-Out</i>	<i>Protocol Addr.</i>	<i>Hardware Addr.</i>
R	5		900	180.3.6.1	ACAE32457342
P	2	2		129.34.4.8	
P	14	5		201.11.56.7	
R	8		450	114.5.7.89	457342ACAE32
P	12	1		220.55.5.7	
P	23	1		116.1.7.22	
R	9		60	19.1.7.82	4573E3242ACA
R	18		900	188.11.8.71	E34573242ACA



EXAMPLE 5

Twenty-five seconds later, the cache-control module updates every entry. The time-out values for the first three resolved entries are decremented by 60. The time-out value for the last resolved entry is decremented by 25. The state of the next-to-the last entry is changed to FREE because the time-out is zero. For each of the three pending entries, the value of the attempts

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Table 7.4 Updated cache table for Example 5

<i>State</i>	<i>Queue</i>	<i>Attempt</i>	<i>Time-Out</i>	<i>Protocol Addr.</i>	<i>Hardware Addr.</i>
R	5		840	180.3.6.1	ACAE32457342
P	2	3		129.34.4.8	
F					
R	8		390	114.5.7.89	457342ACAE32
P	12	2		220.55.5.7	
P	23	2		116.1.7.22	
F					
R	18		875	188.11.8.71	E34573242ACA

7.3 RARP

RARP finds the logical address for a machine that only knows its physical address.

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The topics discussed in this section include:

Packet Format

Encapsulation

RARP Server

Alternative Solutions to RARP

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Note:

*The RARP request packets are broadcast;
the RARP reply packets are unicast.*

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Figure 7.10 *RARP operation*

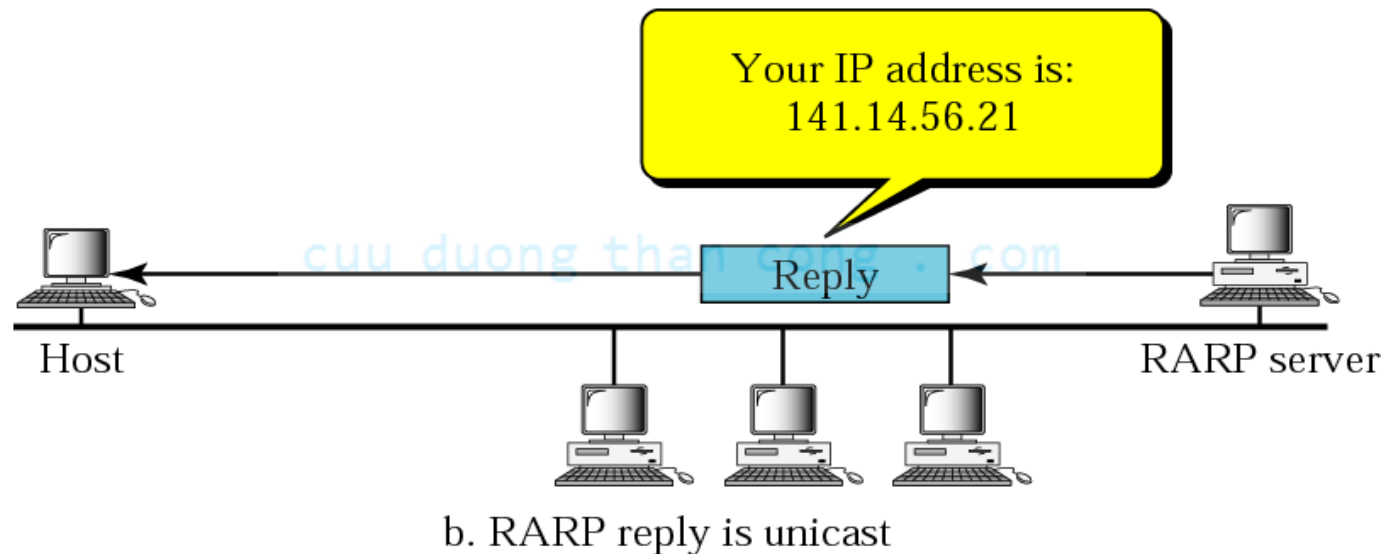
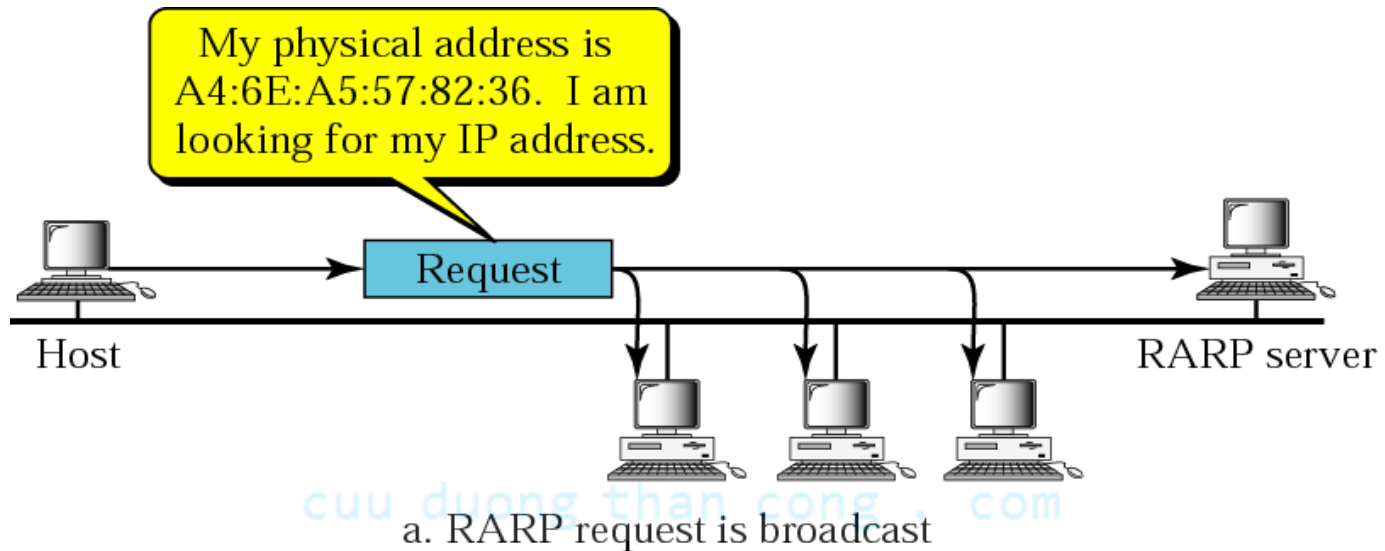


Figure 7.11 *RARP packet*

Hardware type		Protocol type
Hardware length	Protocol length	Operation Request 3, Reply 4
Sender hardware address (For example, 6 bytes for Ethernet)		
Sender protocol address (For example, 4 bytes for IP) (It is not filled for request)		
Target hardware address (For example, 6 bytes for Ethernet) (It is not filled for request)		
Target protocol address (For example, 4 bytes for IP) (It is not filled for request)		

Figure 7.12 *Encapsulation of RARP packet*

