



# Cryptography and Network Security

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## *Chapter 5*

# cuu duong than cong . com Digital Signatures

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# Outline

- Digital Signatures
- Digital Signature Algorithm and Standard

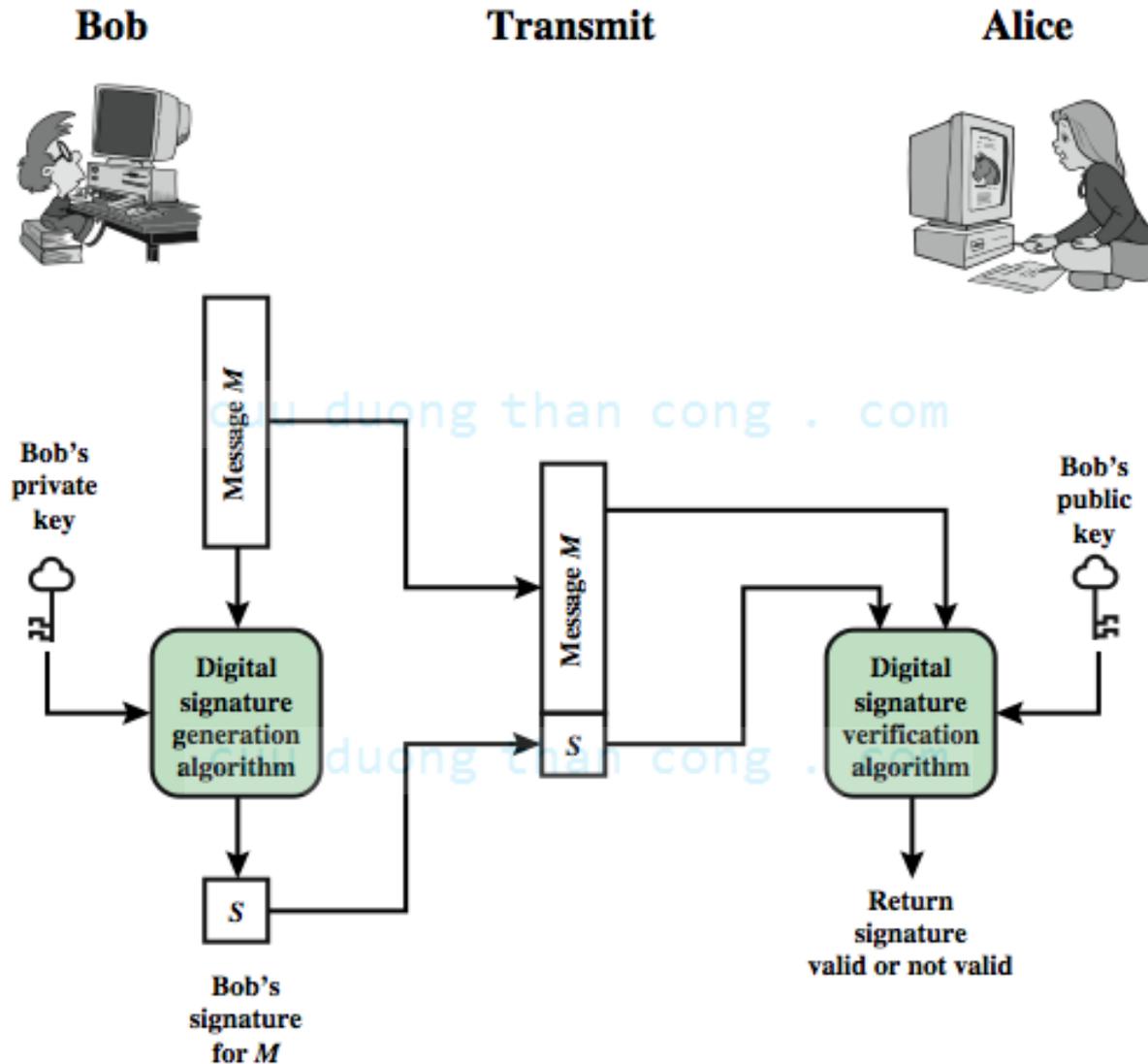
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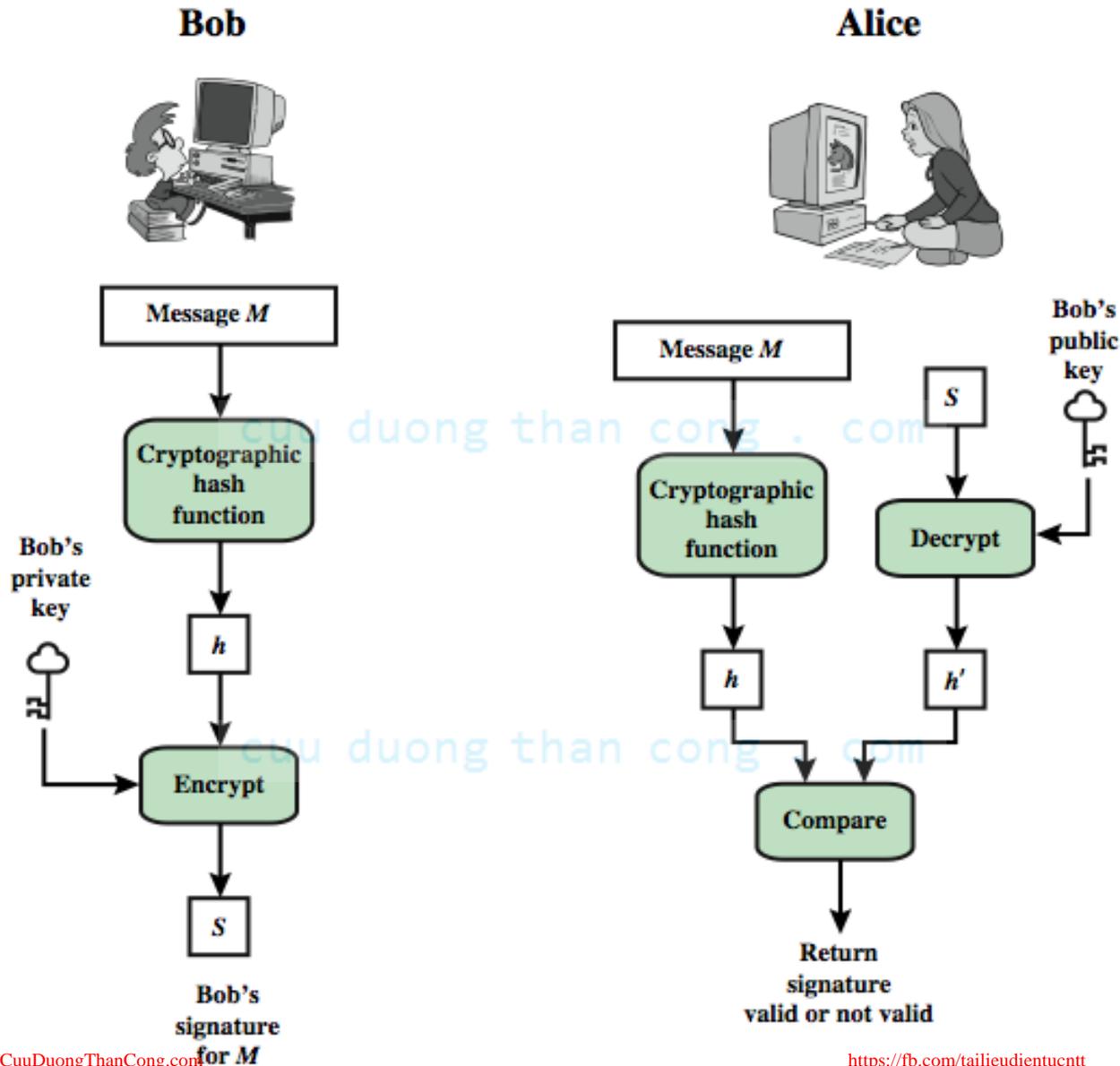
# Digital Signatures

- A **digital signature** is an **authentication mechanism** that enables the creator of a message to attach a code that acts as a signature.
- Typically the signature is formed by taking the **hash of the message** and **encrypting the message with the creator's private key**.
- The signature guarantees the **source** and **integrity** of the message.
- The digital signature standard (DSS) is an NIST standard that uses the **secure hash algorithm** (SHA).

# Digital Signature Model



# Digital Signature Model



# Attacks and Forgeries

## ■ attacks

- key-only attack
- known message attack
- generic chosen message attack
- directed chosen message attack
- adaptive chosen message attack

## ■ break success levels

- total break
- selective forgery
- existential forgery

# Digital Signature Requirements

- must depend on the message signed
- must use information unique to sender
  - to prevent both forgery and denial
- must be relatively easy to produce
- must be relatively easy to recognize & verify
- be computationally infeasible to forge
  - with new message for existing digital signature
  - with fraudulent digital signature for given message
- be practical save digital signature in storage

# Direct Digital Signatures

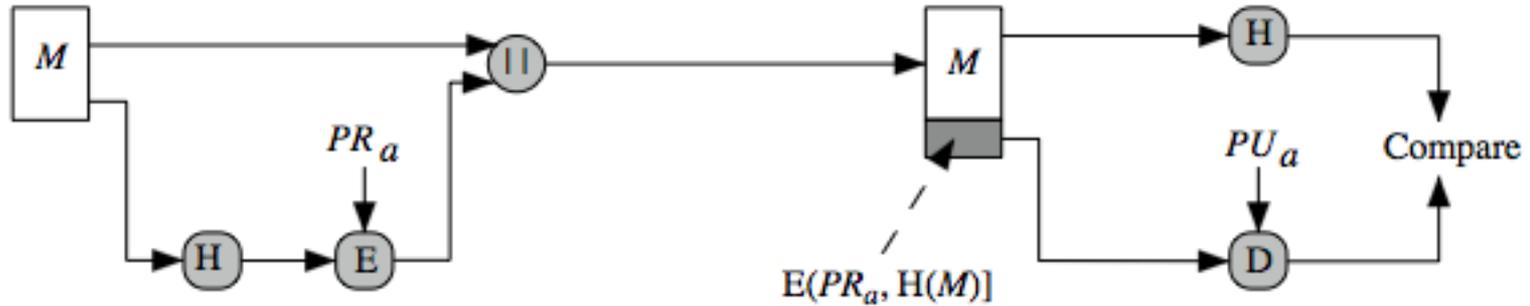
- involve only sender & receiver
- assumed receiver has sender's public-key
- digital signature made by sender signing entire message or hash with private-key
- can encrypt using receivers public-key
- important that sign first then encrypt message & signature
- security depends on sender's private-key

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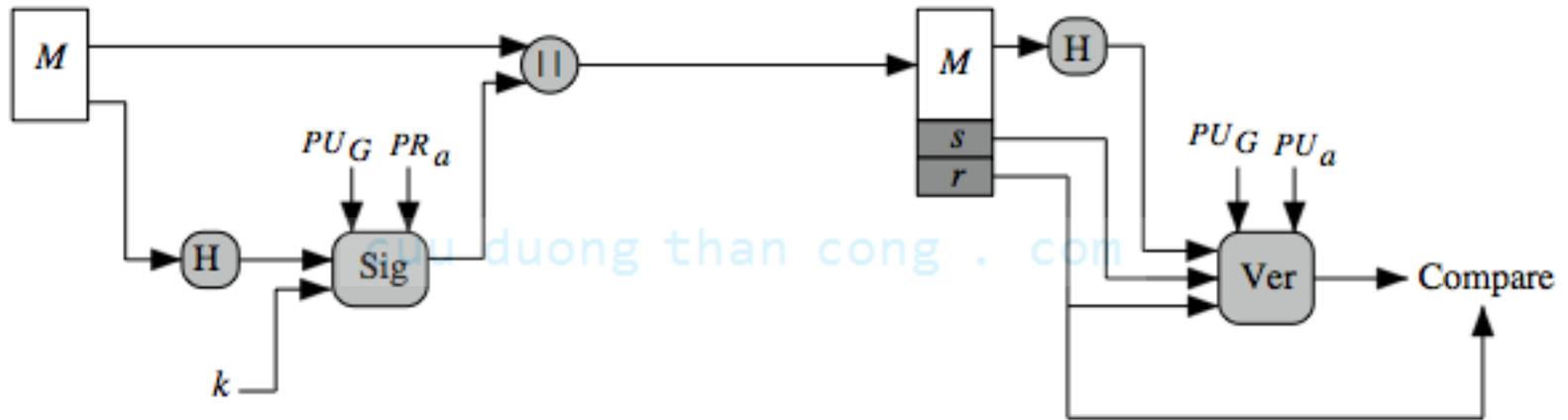
# Digital Signature Standard (DSS)

- US Govt approved signature scheme
- designed by NIST & NSA in early 90's
- published as FIPS-186 in 1991
- revised in 1993, 1996 & then 2000
- uses the SHA hash algorithm
- DSS is the standard, DSA is the algorithm
- FIPS 186-2 (2000) includes alternative RSA & elliptic curve signature variants
- DSA is digital signature only unlike RSA
- is a public-key technique

# DSS vs. RSA Signatures



(a) RSA Approach



(b) DSS Approach

# Digital Signature Algorithm (DSA)

- creates a 320 bit signature
- with 512-1024 bit security
- smaller and faster than RSA
- a digital signature scheme only
- security depends on difficulty of computing discrete logarithms
- variant of ElGamal & Schnorr schemes

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# DSA Key Generation

- **have shared global public key values (p,q,g):**
  - choose 160-bit prime number  $q$
  - choose a large prime  $p$  with  $2^{L-1} < p < 2^L$ 
    - o where  $L = 512$  to  $1024$  bits and is a multiple of  $64$
    - o such that  $q$  is a 160 bit prime divisor of  $(p-1)$
  - choose  $g = h^{(p-1)/q}$ 
    - o where  $1 < h < p-1$  and  $h^{(p-1)/q} \bmod p > 1$
- **users choose private & compute public key:**
  - choose random private key:  $x < q$
  - compute public key:  $y = g^x \bmod p$

# DSA Signature Creation

- **to sign a message  $M$  the sender:**
  - generates a random signature key  $k$ ,  $k < q$
  - Note:  $k$  must be random, be destroyed after use, and never be reused

- **then computes signature pair:**

$$r = (g^k \bmod p) \bmod q$$

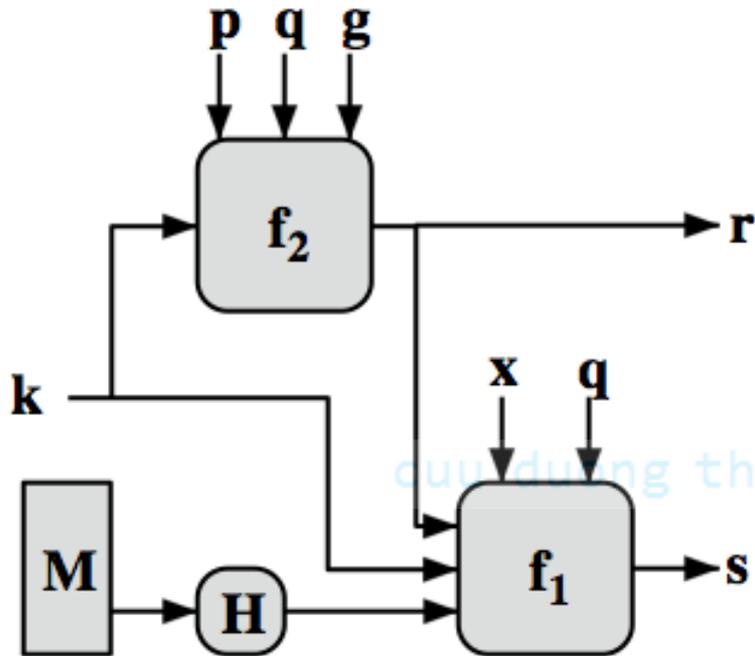
$$s = [k^{-1}(H(M) + xr)] \bmod q$$

- **sends signature  $(r,s)$  with message  $M$**

# DSA Signature Verification

- having received  $M$  & signature  $(r,s)$
- to verify a signature, recipient computes:  
 $w = s^{-1} \bmod q$   
 $u_1 = [H(M)w] \bmod q$   
 $u_2 = (rw) \bmod q$   
 $v = [(g^{u_1} y^{u_2}) \bmod p] \bmod q$
- if  $v=r$  then signature is verified
- see Appendix A for details of proof why

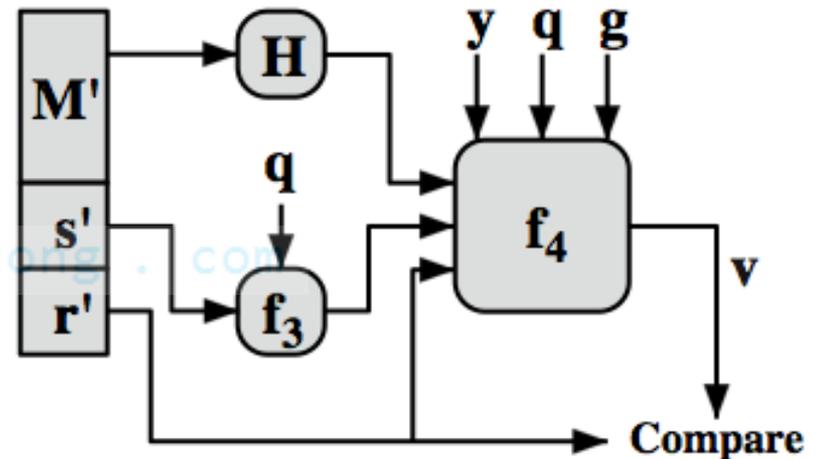
# DSS Overview



$$s = f_1(H(M), k, x, r, q) = (k^{-1} (H(M) + xr)) \bmod q$$

$$r = f_2(k, p, q, g) = (g^k \bmod p) \bmod q$$

(a) Signing



$$w = f_3(s', q) = (s')^{-1} \bmod q$$

$$v = f_4(y, q, g, H(M'), w, r')$$

$$= ((g^{H(M')w} \bmod q \cdot y^{r'w} \bmod q) \bmod p) \bmod q$$

(b) Verifying

# Summary

We have discussed:

- Digital Signatures
- Digital Signature Algorithm and Standard

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# References

1. Cryptography and Network Security, Principles and Practice, William Stallings, Prentice Hall, Fifth Edition, 2011

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